British Informatics Olympiad Final

3–5 April, 1998

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Making a Difference

The UNIX utility `diff` compares two files, and outputs a list of changes necessary for transforming the first into the second. For this question we will look at a possible way of implementing a difference program; one that works with files containing one word per line.

There are three types of transformation : appending, deleting and changing. An output of the form "*n1* `a` *n2*,*n3*" means we append lines *n2* to *n3* of the second file, after line *n1* of the first file. The output "*n1*,*n2* `c` *n3*,*n4*" means change lines *n1* to *n2* of the first file to lines *n3* to *n4* of the second. Finally, "*n1*,*n2* `d` *n3*" tells us to delete lines *n1* to *n2* in the first file. Each of these instructions are followed by the lines that are affected in each file; a '<' indicating a line from the first file, and a '>' indicating a line from the second.

For example, if we compare the following two files

```
she                 the
sells               sea
sea                 shells
shells              she
on                  sells
the                 are
sea                 sea
shore               shells
                    I'm
                    sure
```

a possible response would be :

```
0a1,3
> the
> sea
> shells
2a6,6
> are
5,8c9,10
< on
< the
< sea
< shore
---
> I'm
> sure
```

Question 1.1

Calculate valid difference outputs for the following pairs of files (if written with one word on each line) :

(a) it was the best of times it was the worst of times
it was the age of wisdom it was the age of foolishness

(b) permuting this sentence may render it meaningless
may it sentence permuting render this meaningless

Question 1.2

There may be many ways of transforming one file into another. The method we choose might depend on several criterion; for example we may wish to minimize the number of lines that are transformed, or try to keep large blocks of lines intact. Suppose we are are transforming an n line file into a (different) m line one. What is the minimum number of transformations we can use? What is the maximum number?

Question 1.3

Suppose we have run our program and calculated how to transform *file1* into *file2*. How can we use this information to generate a transformation from *file2* to *file1*, and hence why does the deletion transformation have the format "*n1,n2* `d` *n3*"?

2. Longest common subsequence

A *subsequence* of a sequence X is just X with some of the elements, possibly none, removed. Suppose we have two sequences X and Y; Z is a *common subsequence* of X and Y if it is a subsequence of both of them. If there does not exist a common subsequence with more elements in it, we say Z is a *longest common subsequence* (LCS). [Note that there may be more than one.] Given X and Y we may be interested in determining a longest common subsequence.

Question 2.1

(a) What is the largest common subsequence of $< B, I, O, B, I, O, B, I, O >$ and $< I, O, I, B, I, O, I, O, I >$.

(b) Give a pair a sequences which have more than one LCS, along with example LCSs.

Question 2.2

A naive algorithm for calculating an LCS would be to calculate all the subsequences of X, and check to see which are also subsequences of Y. Why would this impractical for all but the smallest sequences?

We can recursively calculate an LCS of X and Y, by considering the LCSs of pairs of prefixes of X and Y. Let X_n be the sequence $< x_1, x_2, \ldots, x_n >$ and Y_m the sequence $< y_1, y_2, \ldots, y_m >$. Furthermore let $Z_k =< z_1, \ldots, z_k >$ be any LCS of X_n and Y_m.

- If $x_n = y_m$, then $z_k = x_n = y_m$ and Z_{k-1} is an LCS of X_{n-1} and Y_{m-1}.
- If $x_n \neq y_m$, then if $z_k \neq x_n$, Z_k must be an LCS of X_{n-1} and Y_m.
- If $x_n \neq y_m$, then if $z_k \neq y_m$, Z_k must be an LCS of X_n and Y_{m-1}.

Question 2.3

If $x_n \neq y_m$ how does the length of the LCS of X_n and Y_m relate to the LCS of X_{n-1} and Y_m and the LCS of X_n and Y_{m-1}.

Question 2.4

Suppose the X and Y are stored in the arrays `X` and `Y`. Using the above relationships write a recursive procedure `LCS_Length` which is called with the parameters `n` and `m`, and returns the length of the LCS of X_n and Y_m. [Indicate whether your arrays begin at 0 or 1.]

The recursive procedure needs to call itself since we base our solution of the problem, on the results of smaller problems. The flaw with this approach is that `LCS_Length` may get called many times for the same values of `n` and `m`. For example

`LCS_Length(5,5)` may need to look at `LCS_Length(4,5)` and `LCS_Length(5,4)`, both of which may need to look at `LCS_Length(4,4)`.

One way to prevent repeated calculation of the same problem would be to store the results of the sub-problems we have solved. When we are asked to solve a problem for the second time, all we need to do is recall this stored value. The *dynamic programming* approach is to calculate the sub-problems in a different order, so that whenever we need the result of another problem we have already solved it.

The outline of a dynamic programming procedure, might look as follows :

```
for i := 0 to n do
  for j := 0 to m do
    lcsl[i,j] := "length of LCS of X_i and Y_j"
```

This procedure builds up a table `lcsl`, where `lcsl[i,j]` contains the length of the LCS of X_i and X_j. Observe that by the time we calculate `lcsl[i,j]` all the smaller sub-problems have already been solved.

Question 2.5

Suppose we have calculated the table `lcsl`, and hence know the length of the LCS. How can we use this information to find an actual LCS?

3. Creating a difference output

We will now return to the problem of finding a transformation between two files. If we treat the individual words in the two files as the elements of our sequences, we can use our LCS algorithm to find sequences of words in common between the two files. The output from our algorithm might look something like

$$(i_1, j_1), (i_2, j_2), \ldots, (i_k, j_k)$$

where the i_n^{th} word in the first file is the same as the j_n^{th} word in the second, and both $i_1, \ldots, i_k$ and $j_1, \ldots, j_k$ are strictly increasing sequences. For example, if we look back at original example, the output would be

$$(1, 4), (2, 5), (3, 7), (4, 8)$$

Of course, we are really interested in the parts of the files that are *not* common. ie.

- words 1 to $i_1 - 1$ in the first file, compared with 1 to $j_1 - 1$ in the second.
- words $i_1 + 1$ to $i_2 - 1$ in the first file, compared with $j_1 + 1$ to $j_2 - 1$ in the second.

  $\vdots$

- word $i_k + 1$ to the end of the first file, compared with $j_k + 1$ to the end of the second file.

[**NB** "words x to y" where $x > y$, signifies an empty sequence.]

Question 3

Suppose you have (from above) that "words *n1* to *n2* in the first file have nothing in common with words *n3* to *n4* in the second file." When, and how, does this convert to

(a) an 'append' transformation?
(b) a 'delete' transformation?
(c) a 'change' transformation?